Outline

• Polygon clipping
  – Sutherland-Hodgman,
  – Weiler-Atherton

• Polygon filling
  – Scan filling polygons
  – Flood filling polygons

• Introduction and discussion of homework #3
Polygon

- Ordered set of vertices (points)
  - Usually counter-clockwise
- Two consecutive vertices define an edge
- Left side of edge is inside
- Right side is outside
- Last vertex implicitly connected to first
- In 3D vertices should be co-planar
Polygon Clipping

• Lots of different cases

• Issues
  – Edges of polygon need to be tested against clipping rectangle
  – May need to add new edges
  – Edges discarded or divided
  – Multiple polygons can result from a single polygon
The Sutherland-Hodgman Polygon-Clipping Algorithm

• Divide and Conquer

• Idea:
  – Clip single polygon using single infinite clip edge
  – Repeat 4 times

• Note the generality:
  – 2D convex n-gons can clip arbitrary n-gons
  – 3D convex polyhedra can clip arbitrary polyhedra
Sutherland-Hodgman Algorithm

• Input:
  – \( v_1, v_2, \ldots, v_n \) the vertices defining the polygon
  – Single infinite clip edge w/ inside/outside info

• Output:
  – \( v'_1, v'_2, \ldots, v'_m \), vertices of the clipped polygon

• Do this 4 (or \( n_e \)) times

• Traverse vertices (edges)

• Add vertices one-at-a-time to output polygon
  – Use inside/outside info
  – Edge intersections
Sutherland-Hodgman Algorithm

• Can be done incrementally
• If first point inside add. If outside, don’t add
• Move around polygon from $v_1$ to $v_n$ and back to $v_1$
• Check edge $v_i, v_{i+1}$ wrt the clip edge
• Need $v_i, v_{i+1}$’s inside/outside status
• Add vertex one at a time. There are 4 cases:

![Polygon clipping diagram](image-url)
Sutherland-Hodgman Algorithm

- foreach polygon \( P \) \( P' = P \)
  - foreach clipping edge (there are 4) {
    - Clip polygon \( P' \) to clipping edge
      - foreach edge in polygon \( P' \)
        » Check clipping cases (there are 4)
          » Case 1 : Output \( v_{i+1} \)
          » Case 2 : Output intersection point
          » Case 3 : No output
          » Case 4 : Output intersection point
            \& \( v_{i+1} \) \}

Sutherland-Hodgman Algorithm
Sutherland-Hodgman Algorithm

Animated by Max Peysakhov @ Drexel University
Final Result

{A, B, X, Y, E, Z, W, A}

Note: Edges XY and ZW!
Issues with Sutherland-Hodgman Algorithm

- Clipping a concave polygon
- Can produce two CONNECTED areas
Weiler-Atherton Algorithm

- General clipping algorithm for concave polygons with holes
- Produces multiple polygons (with holes)
- Make linked list data structure
- Traverse to make new polygon(s)
Weiler-Atherton Algorithm

- Given polygons A and B as linked list of vertices (counter-clockwise order)
- Find all edge intersections & place in list
- Insert as “intersection” nodes
- Nodes point to A & B
- Determine in/out status of vertices
Linked List Data Structure

Intersection Nodes
Intersection Special Cases

• If “intersecting” edges are parallel, ignore
• Intersection point is a vertex
  – Vertex of A lies on a vertex or edge of B
  – Edge of A runs through a vertex of B
  – Replace vertex with an intersection node
Weiler-Atherton Algorithm: Union

• Find an “outside” vertex
• Traverse linked list
• At each intersection point switch to other polygon
• Do until return to starting vertex
• If there are unvisited “outside” vertices, go to one and repeat
• All visited vertices and nodes define union’ed polygon
Example: Union

{V1, V2, V3, P0, V8, V4, P3, V0},  {V6, P1, P2}
Weiler-Atherton Algorithm: Intersection

- Start at intersection node
  - If connected to an “inside” vertex, go there
  - Else step to an intersection point
  - If neither, stop
- Traverse linked list
- At each intersection point switch to other polygon and remove intersection point from list
- Do until return to starting intersection point
- If intersection list not empty, pick another one
- All visited vertices and nodes define and’ed polygon
Example: Intersection

\{P1, V7, P0\}, \{P3, V5, P2\}
Boolean Special Cases

If polygons don’t intersect

– Union
  • If one inside the other, return polygon that surrounds the other
  • Else, return both polygons

– Intersection
  • If one inside the other, return polygon inside the other
  • Else, return no polygons
Union’ing Two Simple Convex Polygons (A and B)

- Assume that polygon edges are ordered
- Set $P_0 = A$
- Set $P_1 = B$
- Find a vertex $v_i$ of A outside of B
- Add $v_i$ to Output
- Set current edge $E$ as $(v_i, v_{i+1})$
Union’ing Two Simple Convex Polygons (A and B)

• While ((len(Output) < 2) || (Output.first != Output.last)) {
  – Intersect E with all the edges of P1
  – There can be at most two intersections
  – If there are no intersections
    • Add \( v_{i+1} \) (end vertex of E) to Output
    • E = E.next
Union’ing Two Simple Convex Polygons (A and B)

- Else  // There were 1 or 2 intersections
  - Add intersection point with lowest $t$ value along $E$ to Output, i.e. the closest one
  - Add last vertex of P1’s intersected edge to Output
  - Set $E$ equal to next edge of P1
  - Temp = P1
  - P1 = P0  // Switch to the other polygon
  - P0 = Temp

- }  // End of While loop

- Write Output as the Union’ed polygon
Example
Result
Point P Inside a Polygon?

- Connect P with another point P` that you know is outside polygon
- Intersect segment PP` with polygon edges
- Watch out for vertices!
- If # intersections is even (or 0) → Outside
- If odd → Inside
Filling Primitives: Rectangles, Polygons & Circles

• Two part process
  – Which pixels to fill?
  – What values to fill them with?

• Idea: Coherence
  – Spatial: pixels are the same from pixel-to-pixel and scan-line to scan line;
  – Span: all pixels on a span get the same value
  – Scan-line: consecutive scan lines are the same
  – Edge: pixels are the same along edges
Scan Filling Primitives: Rectangles

- Easy algorithm
  - Fill from $x_{\text{min}}$ to $x_{\text{max}}$
    Fill from $y_{\text{min}}$ to $y_{\text{max}}$

- Issues
  - What if two adjacent rectangles share an edge?
  - Color the boundary pixels twice?
  - Rules:
    - Color only interior pixels
    - Color left and bottom edges
Scan Filling Primitives: Polygons

• Observe:
  – FA, DC intersections are integer
  – FE, ED intersections are not integer

• For each scan line, how to figure out which pixels are inside the polygon?
Scan Filling Polygons

- Idea #1: use midpoint algo on each edge, fill in between extrema points
- Note: many extrema pixels lie outside the polygon
- Why: midpoint algo has no sense of in/out

(a)

- Span extrema
- Other pixels in the span
Scan Filling Polygons

- Idea #2: draw pixels only strictly inside
  - Find intersections of scan line with edges
  - Sort intersections by increasing x coordinate
  - Fill pixels on inside based on a parity bit
    - $B_p$ initially even (off)
    - Invert at each intersect
    - Draw when odd, do not draw when even
Scan Filling Polygons

• Issues with Idea #2:
  – If at a fractional x value, how to pick which pixels are in interior?
  – Intersections at integer vertex coordinates?
  – Shared vertices?
  – Vertices that define a horizontal edge?
How to handle vertices?

• Problem:
  – vertices are counted twice

• Solution:
  – If both neighboring vertices are on the same side of the scan line, don’t count it
  – If both neighboring vertices are on different sides of a scan line, count it once
  – Compare current y value with y value of neighboring vertices
Scan-Filling a Polygon
How to handle horizontal edges?

- Idea: don’t count their vertices
- Apply open and closed status to vertices to other edges
  - $y_{\text{min}}$ vertex closed
  - $y_{\text{max}}$ vertex is open
- On AB, A is at $y_{\text{min}}$ for JA; AB does not contribute, $B_p$ is odd and draw AB
- Edge BC has $y_{\text{min}}$ at B, but AB does not contribute, $B_p$ becomes even and drawing stops
How to handle horizontal edges?

- Start drawing at IJ ($B_p$ becomes odd).
- C is $y_{max}$ (open) for BC. $B_p$ doesn’t change.
- Ignore CD. D is $y_{min}$ (closed) for DE. $B_p$ becomes even. Stop drawing.
- I is $y_{max}$ (open) for IJ. No drawing.
- Ignore IH. H is $y_{min}$ (closed) for GH. $B_p$ becomes odd. Draw to FE.
- Ignore GF. No drawing.
Polygon Filling Algorithm

• For each polygon
  – For each edge, mark each scan-line that the edge crosses by examining its $y_{min}$ and $y_{max}$
    • If edge is horizontal, ignore it
    • If $y_{max}$ on scan-line, ignore it
    • If $y_{min} \leq y < y_{max}$ add edge to scan-line $y$‘s edge list
  – For each scan-line between polygon’s $y_{min}$ and $y_{max}$
    • Calculate intersections with edges on list
    • Sort intersections in $x$
    • Perform parity-bit scan-line filling
    • Check for double intersection special case
  – Clear scan-lines’ edge list
How to handle slivers?

- When the scan area does not have an "interior"
- Solution: use anti-aliasing
- But, to do so will require softening the rules about drawing only interior pixels
Scan Filling Curved Objects

- Hard in general case
- Easier for circles and ellipses.
- Use midpoint Alg to generate boundary points.
- Fill in horizontal pixel spans
- Use symmetry
Boundary-Fill Algorithm

- Start with some internal point \((x, y)\)
- Color it
- Check neighbors for filled or border color
- Color neighbors if OK
- Continue recursively
4 Connected Boundary-Fill Alg

Void BoundaryFill4( int x, int y, int fill, int bnd)
{
    If Color(x,y) != fill and Color(x,y) != bnd
    {
        SetColor(x,y) = fill;
        BoundaryFill4(x+1, y, fill, bnd);
        BoundaryFill4(x, y +1, fill, bnd);
        BoundaryFill4(x-1, y, fill, bnd);
        BoundaryFill4(x, y -1, fill, bnd);
    }
}
Boundary-Fill Algorithm

• Issues with recursive boundary-fill algorithm:
  – May make mistakes if parts of the space already filled with the Fill color
  – Requires very big stack size

• More efficient algorithms
  – First color contiguous span along one scan line
  – Only stack beginning positions of neighboring scan lines
Programming Assignment 3

• Read in 2 polygons in Postscript format
• Use simplified Weiler-Atherton algorithm to compute their union
• Write out a single polygon in Postscript format